

# IE170: Algorithms in Systems Engineering: Lecture 29

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## Taking Stock

### Last Time

- Matrix Review

### This Time

- Solving Triangular Systems
- Solving Symmetric Positive Definite Systems
- Least Squares

## Systems of Equations: $Ax = b$

- From our previous discussion, we know that the system of equations  $Ax = b$  has a unique solution if and only if the matrix  $A$  is square and invertible
- This is true if the columns  $A_j$  are linearly independent
- From now on, we will consider only invertible systems.
- In fact, today we will consider special versions of  $A$

### The \$64 Question

- How do we solve a systems of equations?
- We **factor** the matrix  $A$  into a simpler form



## Triangular Systems

- Let's suppose that we are able to find two  $n \times n$  matrices  $L, U$  such that

$$A = LU$$

where

- $L$  is upper triangular.
- $U$  is lower triangular with 1's on the diagonal.
- How could use such a decomposition to solve the system  $Ax = b$ ?

## Using a Triangular Decomposition

- Once we have an triangular decomposition, we can use it to easily solve the system  $Ax = b$ .
- Note that the system  $Ax = b$  is equivalent to the original system, which is then equivalent to  $LUx = b$ .
- We can solve the system in two steps:
  - First solve the system  $Ly = b$  (forward substitution).
  - Then solve the system  $Ux = y$  (backward substitution).



## Forward Substitution

$$\begin{aligned} \ell_{11}y_1 &= b_1 \\ \ell_{21}y_1 + \ell_{22}y_2 + \dots &= b_2 \\ \ell_{n1}y_1 + \ell_{n2}y_2 + \ell_{n3}y_3 + \dots &= b_n \end{aligned}$$

- Just substitute forward into:

$$y_i = \frac{b_i - \sum_{j=1}^{i-1} \ell_{ij}y_j}{\ell_{ii}}$$

- So we have  $y$  such that  $Ly = b$ .



## Example Matrix

- Next, we simply solve the system  $Ux = y$
- Backwards substitution works in a similar fashion, but loops “down” from  $n$  down to 1

$$\begin{aligned} u_{11}x_1 + u_{12}x_2 + u_{13}x_3 + \dots &= y_1 \\ &u_{22}x_2 + u_{23}x_3 + \dots = y_2 \\ &&u_{nn}x_n = y_n \end{aligned}$$

TRIANGULARSOLVE( $L, U, b$ )

```
1  $n \leftarrow \text{rows}[L]$ 
2 for  $i \leftarrow 1$  to  $n$ 
3 do  $y[i] \leftarrow (b[i] - \sum_{j<i} \ell_{ij}y_j) / \ell_{ii}$ 
4 for  $i \leftarrow n$  to 1
5 do  $x[i] \leftarrow (y[i] - \sum_{j>i} u_{ij}x_j) / u_{ii}$ 
6 return  $x$ 
```



## Example

$$L = \begin{bmatrix} 4 & 0 & 0 & 0 \\ 1 & 3 & 0 & 0 \\ 2 & 2 & 2 & 0 \\ 1 & 1 & 3 & 1 \end{bmatrix} \quad U = L^T \quad b = \begin{bmatrix} 32 \\ 26 \\ 28 \\ 30 \end{bmatrix}$$

- We'll solve  $Ax = LUx = b$ .
- Hopefully we get  $x = (1, 1, 1, 1)^T$



## Special Matrices

- If  $A$  is a square symmetric ( $A = A^T$ ) matrix such that

$$x^T Ax > 0 \quad \forall x \in \mathbb{R}^n, x \neq 0$$

then  $A$  is said to be **symmetric positive definite**.

### They are everywhere!

- Electrical circuit problems
- Structural Engineering (elastic deformations)
- Variance-Covariance matrices
- Numerical solution of partial differential equations
- Solution of linear systems
  - $Ax = b \Leftrightarrow x$  minimizes  $1/2x^T Ax - b^T x$
- Least Squares Problems!**



## Fun SPD Facts

- If  $A$  is SPD, then  $A$  is nonsingular
  - Proof:** If  $A$  is singular, then there is  $x \neq 0 \in \mathbb{R}^n$  with  $Ax = 0$ , so  $x^T Ax = 0$ , and  $x$  is not SPD. QUITE ENOUGH DONE
- If  $M$  is non-singular, then  $A = MM^T$  is SPD
  - Proof:**  $A^T = (MM^T)^T = MM^T = A$ , so  $A$  is symmetric.  $x^T Ax = x^T (MM^T)x = y^T y$  for  $y = M^T x$ .  $y^T y = \sum_{i=1}^n y_i^2 \geq 0$ , and it is only 0 when  $y = 0$ , but  $y \neq 0$  or else  $M^T$  would have been singular, since  $y = M^T x$ . QUITE ENOUGH DONE



## Looking for a Decomposition

- Wouldn't it be awesome if  $U = L^T$  (like our example), so there was a decomposition of the form  $A = LL^T$ . Let's check to see if that is possible:

$$\begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix} = \begin{bmatrix} \ell_{11} & 0 & \cdots & 0 \\ \ell_{21} & \ell_{22} & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots \\ \ell_{n1} & \ell_{n2} & \cdots & \ell_{nn} \end{bmatrix} \begin{bmatrix} \ell_{11} & \ell_{21} & \cdots & \ell_{n1} \\ 0 & \ell_{22} & \cdots & \ell_{n2} \\ \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & \ell_{nn} \end{bmatrix}$$

- And this implies that...

- $a_{i1} = \ell_{i1}\ell_{11}$ , so  $a_{11} = \ell_{11}^2$ ,  $\ell_{i1} = a_{i1}/\ell_{11}$
- $a_{i2} = \ell_{i1}\ell_{21} + \ell_{i2}\ell_{22}$



## General Formula

### In General

$$a_{ij} = g_{i1}g_{j1} + g_{i2}g_{j2} + \cdots + g_{i,j-1}g_{j,j-1} + g_{ij}g_{jj}$$

- This **only** depends on columns up to  $j$ !
- Assuming we have computed the first  $j - 1$  columns of  $L$ , the  $j^{\text{th}}$  columns can be computed using the formulae

$$g_{jj} = \sqrt{a_{jj} - \sum_{k < j} g_{jk}^2}$$

$$g_{ij} = \frac{a_{ij} - \sum_{k < j} g_{ik}g_{jk}}{g_{jj}} \quad \text{for } j > i$$



## Query and Example

- What if  $a_{jj} - \sum_{k<j} g_{jk}^2 < 0$ ?
- Then  $A$  is **not** SPD.
- The proof of this fact is too complicated to give now, but it **is** true that  $A$  is SPD **if and only if** it can be written as  $A = LL^T$  for a lower triangular matrix  $L$
- $L$  is known as the **Cholesky** factor of  $A$ , after the French mathematician André-Louis Cholesky.

$$A = \begin{bmatrix} 16 & 4 & 8 & 4 \\ 4 & 10 & 8 & 4 \\ 8 & 8 & 12 & 10 \\ 4 & 4 & 10 & 12 \end{bmatrix}$$

- Assuming we can do the arithmetic correctly, we should get  $A = LL^T$ , with  $L$  the previous  $L$  in this lecture.



## Least Squares

- Suppose I am given some data points (measurements)

$$(x_1, y_1), (x_2, y_2), \dots, (x_m, y_m)$$

And we wish to find a function that closely approximates these measurements:

$$y_i = F(x_i) + \eta_i$$

where the  $\eta_i$  are “small.”

- We will assume that  $F(x)$  has the form:

$$F(x) = \sum_{j=1}^n c_j f_j(x)$$



## Least Squares

- A common choice of the “basis functions”  $f_j(x)$  are small order polynomials:

$$F(x) = c_1 + c_2x + c_3x^2 + \dots + c_nx^{n-1}.$$

- Choosing  $n = m$  means that the function will **exactly** match the  $y_i$ , and is generally a “bad idea”, as this is known as **overfitting**,
- Instead,  $n$  is typically much smaller than  $m$ 
  - For example, if  $n = 2$ , then we are looking for the best “linear” fit of the data



## More Least Squares

- Let's create the matrix

$$A = \begin{bmatrix} f_1(x_1) & f_2(x_1) & \vdots & f_n(x_1) \\ f_1(x_2) & f_2(x_2) & \vdots & f_n(x_2) \\ \vdots & \vdots & \dots & \vdots \\ f_n(x_1) & f_n(x_1) & \vdots & f_n(x_1) \end{bmatrix}$$

- So  $Ac = [F(x_1), F(x_2), \dots, F(x_m)]^T$  is the  $m$ -vector of **predicted** values for  $y$ , so
- $\eta = Ac - y$  is the vector that we are trying to minimize



## Least Squares

- In least-squares, we minimize the squared (Euclidean) length of  $\eta$ , or

$$\min \|\eta\|^2 = \|Ac - y\|^2 = \sum_{i=1}^m \left( \sum_{j=1}^n a_{ij}c_j - y_i \right)^2$$

- Taking derivatives, setting the result equal to zero, and putting things back in matrix notation, means that we look for a  $c$  such that

$$(Ac - y)^T A = 0 \text{ or } A^T Ac = A^T y.$$



## Solving Least Squares

### Normal Equations

$$A^T Ac = A^T y$$

- We seek  $c = (A^T A)^{-1} A^T y$
- Sometimes  $A^+ \stackrel{\text{def}}{=} (A^T A)^{-1} A^T$  is call the **pseudoinverse** of  $A$ , and it exists for **non-square**  $A$
- We don't **really** need to take the inverse, we just solve the **SPD system**  $A^T Ac = A^T y$
- This is what you get to do in lab today!

