Solving Linear Systems

IE170: Algorithms in Systems Engineering: Lecture 30

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- Last time we learned about how to solve systems Ax = b, when A was symmetric and positive-definite.
- $\bullet\,$ The key was to factor the matrix into two triangular matrices A=LU
- In the case that A is ${\rm SPD},$ then we can always do this, and in fact $U=L^T.$
- What if A is not SPD?
- The workhorse in this case is the LU-decomposition
- LU-decomposition is very related to (the well-known) Gaussian elimination, a fact we will try to make clear today...





Gaussian Elimination

• An example for today. Let's solve it...

- Subtract twice first equation from the second
- Subtract 3 times the first equation from the third
- Then add 4 times second equation to the third
- You've made a triangular system!
- What were the matrices that produce this?

Elementary Dear Watson!



- We reduced the columns by taking linear combinations of the rows of the matrix.
- This implies that the reduction process can be thought of as a multiplication of A on the left by some matrix
- What does the matrix look like?
- It is an elementary matrix of the form

$$E = I - uv^T$$

 In fact, it's a special form of an elementary matrix: It will be a unit lower triangular matrix with multipliers only in one column



The Elimination Matrix

• Let's find a matrix M_1 that reduces the first column of A.

$$M_1 = \begin{pmatrix} 1 \\ -m_{21} & 1 \\ -m_{31} & 1 & 1 \\ \vdots & \vdots & \ddots \\ -m_{n1} & 0 & 0 & \cdots & 1 \end{pmatrix}$$

- By our properties of matrix multiplication, this matrix
 - Leaves the first row of A alone
 - $\bullet~{\rm Takes}~-m_{21}$ times the first row, adds the second row, and puts this in the second row of the new matrix M_1A
 - Takes $-m_{31}$ times the first row of A, adds the third row, and puts this in the third row of the new matrix M_1A
 - (And So On...)

The UpShot

• To eliminate the first column, we want

$$m_{i1} = \frac{a_{i1}}{a_{11}}$$

- Note: These were exactly the multipliers we used in our simple example
- a_{11} is called the pivot element, and this reduction only works if the pivot element is $\neq 0$
- Next time: What happens if pivot element is 0 (or small)





Let's Carry On

$$M_1 A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ 0 & a_{22}^{(2)} & \cdots & a_{2n}^{(2)} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & a_{n2}^{(2)} & \cdots & a_{nn}^{(2)} \end{pmatrix}$$

• To reduce the second column, we would like a matrix



Matrix Effect

- Again, the matrix M_2 will...
 - Leave first row of $M_1 A$ unchanged in $M_2 M_1 A$
 - $\bullet\,$ Leave second row of M_1A unchanged M_2M_1A
 - Take $-m_{32}$ times second row + third row in M_2M_1A

Lather Rinse Repeat

- Repeat this n-1 times
- In the end, we get $M_{n-1} \cdots M_2 M_1 A = U$
- Fact: The product of unit lower triangular matrices is unit lower triangular
- So in the end we have MA = U, with M unit lower triangular, and U upper traiangular
- This process is known as Gaussian Elimination, and the matrix *M* is known as the product form of the LU factorization

Finding LU Directly

- Here is a recursive method for finding the LU factorization
- We'll divide the matrix A into four pieces:

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix}$$
(1)
$$= \begin{bmatrix} a_{11} & w^{T} \\ v & A' \end{bmatrix}$$
(2)

• Next, we'll use use row operations to change v into the zero vector and record the operations in another matrix.

Finding the LU Decomposition (cont.)

• By simple multiplication, you can verify the following factorization of *A*:

$$A = \begin{bmatrix} a_{11} & w^T \\ v & A' \end{bmatrix}$$
(3)

$$= \begin{bmatrix} 1 & 0 \\ v/a_{11} & I \end{bmatrix} \begin{bmatrix} a_{11} & w^T \\ 0 & A' - vw^T/a_{11} \end{bmatrix}$$
(4)

- We can show that if A is nonsingular, then so is $A' vw^T/a_{11}$.
- So we can recursively call the method to factor the $(n-1) \times (n-1)$ matrix $A' vw^T/a_{11}$.
- \bullet Applying this recursion \boldsymbol{n} times yields the desired factorization



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Finding the LU Decomposition (cont.)

• To see how to get the factorization from the recursive application of the algorithm, we have the following.

$$A = \begin{bmatrix} 1 & 0 \\ v/a_{11} & I \end{bmatrix} \begin{bmatrix} a_{11} & w^T \\ 0 & A' - vw^T/a_{11} \end{bmatrix}$$
(5)

$$= \begin{bmatrix} 1 & 0 \\ v/a_{11} & I \end{bmatrix} \begin{bmatrix} a_{11} & w^T \\ 0 & L'U' \end{bmatrix}$$
(6)

$$= \begin{bmatrix} 1 & 0 \\ v/a_{11} & L' \end{bmatrix} \begin{bmatrix} a_{11} & w^T \\ 0 & U' \end{bmatrix}$$
(7)

- This shows how to obtain the factorization recursively.
- This can also be done iteratively and "in place."

The Algorithm

LU-DECOMPOSITION(
$$A$$
)
1 $n \leftarrow rows[L]$
2 for $k \leftarrow 1$ to n
3 do
4 $u_{kk} \leftarrow a_{kk}$
5 for $i \leftarrow 1$ to n
6 do
7 $\ell_{ik} \leftarrow a_{ik}/u_{kk}$
8 $u_{ki} \leftarrow a_{ki}$
9 for $i \leftarrow k + 1$ to n
10 do
11 for $j \leftarrow k + 1$ to n
12 do
13 $a_{ij} \leftarrow a_{ij} - \ell_{ik}u_{kj}$

$LU \approx Gaussian Elimination$

• We either have A = LU or we have MA = U

- M is unit lower triangular, and in fact the inverse of a unit lower triangular matrix is unit lower-triangular, so $A = M^{-1}U$, and since the elements of L and U are unique, it must be that $L = M^{-1}$
- Because of the special structure of M, we have a (fairly) remarkable relationship

$LU \approx Gaussian Elimination$

• The relationship is the following:

$$M^{-1} = (M_{n-1} \cdots M_2 M_1)^{-1} = M_1^{-1} M_2^{-1} \cdots = L$$
$$L = \begin{pmatrix} 1 & & \\ m_{21} & 1 & & \\ m_{31} & m_{32} & 1 & \\ \vdots & \ddots & \\ m_{n1} & m_{n2} & m_{n3} & \cdots & 1 \end{pmatrix}$$

where the m_{ik} are the multipliers from Gaussian elimination!

• So L and U can be derived directly from the elimination process:

 $a^{(k)}$

(1)

$$\ell_{ik} = m_{ik} = \frac{u_{ik}}{a_{kk}^{(k)}} \qquad u_{kj} = a_{kj}^{(k)}$$

Lecture Notes