

# IE418: Integer Programming

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## Key Things We Learned Last Time

- If  $\{x \in \mathbb{R}^n \mid Ax = b\} \neq \emptyset$ , the maximum number of affinely independent solutions of  $Ax = b$  is  $n + 1 - \text{rank}(A)$ .
- A polyhedron  $P$  is of **dimension  $k$** , denoted  $\text{dim}(P) = k$ , if the maximum number of affinely independent points in  $P$  is  $k + 1$ .
- **2.4** If  $P \subseteq \mathbb{R}^n$ , then  $\text{dim}(P) + \text{rank}(A^=, b^=) = n$
- Partition the inequalities defining  $P$ :
  - $M = \{1, \dots, m\}$ ,
  - $M^= = \{i \in M \mid a_i x = b_i \ \forall x \in P\}$  (the **equality set**),
  - $M^{\leq} = M \setminus M^=$  (the **inequality set**).

## Our Friend POLLY

$$\begin{aligned} x_1 - 2x_2 + x_3 - x_4 + 2x_5 &\leq 3 \\ x_1 - x_5 &\leq 0 \\ -x_1 + x_5 &\leq 0 \\ 2x_2 - x_3 + x_4 &\leq 2 \\ -4x_2 + 2x_3 - 2x_4 &\leq -4 \\ 3x_1 - x_2 &\leq 2 \\ -x_1 &\leq 0 \\ -x_2 &\leq 0 \\ -x_3 &\leq 0 \\ -x_4 &\leq 0 \\ -x_5 &\leq 0 \end{aligned}$$

- $POLLY \subseteq \mathbb{R}^5$ :
- Remember:  
 $\text{dim}(POLLY) = 3$



## Valid Inequalities and Faces

$$P = \{x \in \mathbb{R}^n \mid Ax \leq b\}$$

- The inequality denoted by  $(\pi, \pi_0)$  is called a **valid inequality** for  $P$  if  $\pi x \leq \pi_0 \ \forall x \in P$ .
- Note that  $(\pi, \pi_0)$  is a valid inequality if and only if  $P$  lies in the half-space  $\{x \in \mathbb{R}^n \mid \pi x \leq \pi_0\}$ .
- If  $(\pi, \pi_0)$  is a valid inequality for  $P$  and  $F = \{x \in P \mid \pi x = \pi_0\}$ ,  $F$  is called a **face** of  $P$  and we say that  $(\pi, \pi_0)$  **represents** or **defines**  $F$ .
- A face is said to be **proper** if  $F \neq \emptyset$  and  $F \neq P$ .
- Note that a face has multiple representations.



## More on Faces

- The face represented by  $(\pi, \pi_0)$  is nonempty if and only if  $\max\{\pi x \mid x \in P\} = \pi_0$ .
- If the face  $F$  is nonempty, we say it **supports**  $P$ .
- **3.1** Let  $P$  be a polyhedron with equality set  $M^=$ . If  $F = \{x \in P \mid \pi^T x = \pi_0\}$  is nonempty, then  $F$  is a polyhedron. Let  $M_F^= \supseteq M^=$ ,  $M_F^< = M \setminus M_F^=$ . Then  $F = \{x \mid a_i^T x = b_i \forall i \in M_F^=, a_i^T x \leq b_i \forall i \in M_F^<\}$ 
  - We get the polyhedron  $F$  by taking some of the inequalities of  $P$  and making them equalities
  - The number of distinct faces of  $P$  is finite.



## Facets

- A face  $F$  is said to be a **facet** of  $P$  if  $\dim(F) = \dim(P) - 1$ .
- Facets are all we need to describe polyhedra.
- **3.2** If  $F$  is a facet of  $P$ , then in any description of  $P$ , there exists some inequality representing  $F$ . (By setting the inequality to equality, we get  $F$ ).
- **3.3 and 3.4** Every inequality that represents a face that is not a facet is unnecessary in the description of  $P$ .



## POLLY's Faces

Consider the face

$$F = \{x \in POLLY \mid 2x_1 + 10x_2 - 5x_3 + 5x_4 - 3x_5 = 10\}$$

- Is it proper?
  - $F = POLLY$ ?
  - $F = \emptyset$ ?
- Points:  $(0, 2, 2, 0, 0)$ ,  $(0, 1, 0, 0, 0)$ ,  $(0, 0, 0, 2, 0)$ 
  - Are they in  $F$ ? (Don't forget, they must also be in  $P$ )
  - Are they affinely independent?
  - **Yes!** so  $\dim(F) \geq 2$
  - Is  $\dim(F) \leq 2$ ? (Yes!)
  - $\dim(POLLY) = 3$ , so  $F$  is a facet of  $POLLY$



## Facet Representation

- Remember **3.2**. If  $F$  is a facet of  $POLLY$ , then there is some inequality  $a_k^T x \leq b_k, k \in M^{\leq}$  representing  $F$ .
  - Which inequality in the inequality set of  $POLLY$  represents  $F$ ?
- 
- $x \in POLLY, 2x_1 + 10x_2 - 5x_3 + 5x_4 - 3x_5 = 10 \Rightarrow x_1 = 0$ 
    - So  $F$  is represented by the inequality  $-x_1 \leq 0$



## Polyhedra—A Fundamental Representation Theorem

Putting together what we have seen so far, we can say the following: **(3.5)**

- Every full-dimensional polyhedron  $P$  has a unique (up to scalar multiplication) representation that consists of one inequality representing each facet of  $P$ .
- If  $\dim(P) = n - k$  with  $k > 0$ , then  $P$  is described by a maximal set of linearly independent rows of  $(A^=, b^=)$ , as well as one inequality representing each facet of  $P$ .



## Polyhedra—A Useful Facet Proving Theorem

- Put another way, if a facet  $F$  of  $P$  is represented by  $(\pi, \pi_0)$ , then the set of all representations of  $F$  is obtained by taking scalar multiples of  $(\pi, \pi_0)$  plus linear combinations of the equality set of  $P$ .
- We can use this to actually prove an inequality is a facet! **(3.6)**

- Let  $M^= \equiv (A^=, b^=)$  be the equality set of  $P \subseteq \mathbb{R}^n$ , and let  $F = \{x \in P \mid \pi^T x = \pi_0\}$  be a proper face of  $P$ . The following statements are equivalent

- $F$  is a facet of  $P$
- If  $\lambda x = \lambda_0 \forall x \in F$ , then

$$(\lambda, \lambda_0) = (\alpha\pi + uA^=, \alpha\pi_0 + u^t b^=),$$

for some  $\alpha \in \mathbb{R}, u \in \mathbb{R}^{|M^=|}$



## More Insight—Proving Facets

- This is just an indirect but very useful way to verify affine independence of points.
  - Here we assume that  $P$  is full dimensional  $\dim(P) = n$  (though you can still use Theorem 3.6 even if not).
- Given valid inequality  $\pi^T x \leq \pi_0 \dots$ 
  - 1 Choose  $t \geq n$  points  $x^1, x^2, \dots, x^t$  all satisfying  $\pi^T x = \pi_0$ . Suppose that all these points also lie in a generic hyperplane  $\lambda^T x = \lambda_0$ .

- 2 Solve the linear equation system:

$$\sum_{j=1}^n \lambda_j x_j^k = \lambda_0 \quad \forall k = 1, 2, \dots, t$$

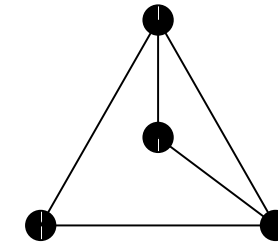
- 3 If the only solution is  $(\lambda, \lambda_0) = \alpha(\pi, \pi_0)$  for  $\alpha \neq 0$ , then  $\pi^T x \leq \pi_0$  is facet defining.



## Example: The Node Packing Polytope

- Given a graph  $G = (V, E)$ , with  $|V| = n$

$$\text{PACK}(G) = \{x \in \mathbb{B}^n \mid x_i + x_j \leq 1 \quad \forall (i, j) \in E\}$$



$$\text{PACK} = \text{conv} \left( \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \\ 0 \\ 1 \end{pmatrix} \right)$$



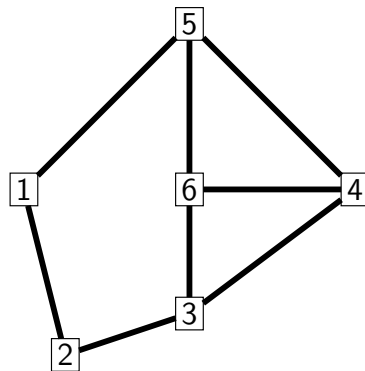
## Example

### Our Task

Prove that

$$x_1 + x_2 + x_3 + x_4 + x_5 \leq 2$$

is a facet-defining inequality for  $\text{PACK}(\hat{G})$ :



## Steps of the Proof

Your writing here

- 1 What is  $\dim(\text{PACK}(\hat{G}))$ ?
- 2 So then  $\text{rank}(A^-, b^-) = 0$
- 3 Let  $F = \{x \in \text{PACK}(\hat{G}) \mid x_1 + x_2 + x_3 + x_4 + x_5 = 2\}$ .  
Suppose  $\lambda^T x = \lambda_0 \forall x \in F$ .  
We will show that  $(\lambda, \lambda_0)$  is a scalar multiple of  $[(1, 1, 1, 1, 1, 0)^T, 2]$ .



## An easier way?

- What we've just done is really just another way to show that the points we chose were affinely independent.

```
octave:1> A = [ 1, 0, 1, 0, 0, 0; 0, 1, 0, 1, 0, 0;
0, 0, 1, 0, 1, 0; 1, 0, 0, 1, 0, 0;
0, 1, 0, 0, 1, 0; 1, 0, 1, 0, 0, 1 ]
```

A =

```
1 0 1 0 0 0
0 1 0 1 0 0
0 0 1 0 1 0
1 0 0 1 0 0
0 1 0 0 1 0
1 0 1 0 0 1
```

```
octave:2> rank(A)
ans = 6
```



## A More Abstract Example

- Let  $C \subseteq V$  be a maximal clique in  $G$ . We will show (two ways) that

$$\sum_{i \in C} x_i \leq 1$$

- is a facet-defining inequality (a facet) of  $\text{PACK}(G)$ .
- First question: What is  $\dim(\text{PACK})$ ?
  - How do we establish the dimension of a polyhedron?



## Way #1—Direct

- To show that  $\sum_{i \in C} x_i \leq 1$  is a facet (that its dimension is  $n - 1$ ), we can give  $n$  affinely independent points in PACK that satisfy  $\sum_{i \in C} x_i = 1$ 
  - Since the hyperplane  $\sum_{i \in C} x_i = 1$  does not contain the origin, this is equivalent to giving  $n$  linearly independent points.
- WLOG, let the clique be  $C = \{1, 2, \dots, k\}$
- Key:  $\forall p \in V \setminus C \exists i_p \in C$  such that  $(i_p, p) \notin E$ . **Why?**
- Points:  $(e_1, e_2, \dots, e_k, e_{k+1} + e_{i_{k+1}}, \dots, e_p + e_{i_p}, \dots, e_n + e_{i_n})$



## Way #2—Indirect

- Let  $F = \{x \in \text{PACK} \mid \sum_{i \in C} x_i = 1\}$
- Suppose  $F \subseteq G \equiv \{x \in \text{PACK} \mid \lambda^T x = \lambda_0\}$  ( $\lambda \neq 0$ )
- If we can show that  $G$  is just a (non-zero) scalar multiple of  $F$ , then we have established that  $F$  is a facet.
- Again, WLOG, let  $C = \{1, 2, \dots, k\}$
- For  $i \leq k$  consider the point  $e_i$ .
  - Satisfies equality  $F$ .  $F \subseteq G \Rightarrow \lambda_i = \lambda_0 \forall i \in C$



## Indirect Facet Proof, cont.

- For  $p \in V \setminus C$ , consider the point  $e_p + e_{i_p}$ : (1's in the coordinates  $p$  and  $i_p$ )
- By the same argument as the previous proof, this point packs, and we can always find such a point  $\forall p \in V \setminus C$
- This point satisfies equality  $F$ .  $F \subseteq G \Rightarrow \lambda_{i_p} + \lambda_p = \lambda_0$
- $\lambda_{i_p} = \lambda_0$ , so  $\lambda_p = \lambda_0 \forall p \in V \setminus C$ .
- So out inequality defining  $G$  looks like  $\lambda_0 \sum_{i \in C} x_i = \lambda_0$ .
- This is a scalar multiple of the inequality defining  $F$ , so  $F$  is a facet defining inequality.
  - $\lambda_0 \neq 0$  since  $\lambda \neq 0$ ,  $\lambda_i = \lambda_0 \forall i \in C$ ,  $\lambda_p = 0 \forall p \in V \setminus C$

