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- Review—Monte Carlo Methods
- "Interior" Sampling Methods
  - ♦ Stochastic Decomposition



$$v^* = \min_{x \in S} \{ f(x) \equiv \mathbb{E}_P g(x; \xi) \}$$

- How can I use sampling to get a (statistical) lower bound on  $v^*$ ?
- How can I use sampling to get a (statictical) upper bound on  $v^*$ ?
- When do you think you will ever get your homeworks back?

#### Putting it all together

- $\widehat{f}_N(x)$  is the sample average function
  - $\diamond \ \mathrm{Draw} \ \omega^1, \dots \omega^N \ \mathrm{from} \ P$
  - $\diamond \ \widehat{f}_N(x) \equiv N^{-1} \sum_{j=1}^N g(x, \omega^j)$
  - $\diamond \ \mathbb{E}[\widehat{f}_N(x)] = f(x) \ge v^*$
- $\hat{v}_N$  is the optimal solution value for the sample average function:

$$\widehat{v}_N \equiv \min_{x \in S} \left\{ \widehat{f}_N(x) := N^{-1} \sum_{j=1}^N g(x, \omega^j) \right\}$$
  
$$\widehat{\mathbb{E}}[\widehat{v}_N] \leq v^*$$

- Estimate  $\mathbb{E}(\hat{v}_N)$  as  $\widehat{\mathbb{E}(\hat{v}_N)} = L_{N,M} = M^{-1} \sum_{j=1}^M \hat{v}_N^j$ 
  - $\diamond\,$  Solve M stochastic LP's, each of sampled size N.



### **Sampling Methods**

- "Interior" sampling methods.
- Sample during the course of the algorithm
  - LShaped Method (Dantzig and Infanger)
  - Stochastic Decomposition (Higle and Sen)
  - Stochastic Quasi-gradient methods (Ermoliev)
- "Exterior" sampling methods
  - ◇ Sample. Then solve problem approximating problem.
  - ♦ Can we get (statistical) bounds on key solution quantities?

- Designed to solve 2-stage stochastic LP with fixed recourse
- Assume randomness only in objective function
- Assume relatively complete recourse and  $Q(x) \ge 0$ .

• 
$$X = \{\Re^n_+ | Ax = b\}$$

$$\min_{x \in X} \mathcal{Q}(x) \equiv \mathbb{E}_P Q(x, \omega) = \int_{\Omega} Q(x, \omega) dP(\omega)$$

$$Q(x,\omega) = \min_{y \in \Re^p_+} \{q^T y : Wy = h(\omega) - T(\omega)x\}$$

# Key Ideas–1

- Do sampling within the optimization
- Take one new sample per iteration
- A bit like "Using Bounds in the LShaped Method"
- There, we used Jensen's Inequality to give a lower bound.
  - There the lower bound was based on using a limited number of samples (namely the mean) and conditioning (to be in a region of Ω to get progressively tigther lower bounds).
- Here, we just randomly draw a sample, and we produce a lower bound that is valid in expectation only.



- Method for "approximately" solving lots of LPs fast
- Cuts "phase out" as they get older. (Since they were based on less sampled information).

#### **Bounds in LShaped Method—Review**

- Use the LShaped method to optimize the problem using Q<sub>L</sub>(x).
  ◊ Only include ω̄ "scenarios".
- When optimized with respect to  $Q_L(x)$ , compare to  $Q_U(x)$ , if we have one
- If  $Q_U(x) Q_L(x)$  is "sufficiently small". Stop.
- Otherwise, refine the partition (which improves the bounds), and repeat.

#### More Detail

Partition form of Jensen's inequality...

• Let  $S = \{\Omega^l, l = 1, 2, \dots v\}$  be some partition of  $\Omega$ :

$$\mathbb{E}_{\omega}[Q(\hat{x},\omega)] \ge \sum_{l=1}^{v} P(\omega \in \Omega^{l})Q(\hat{x},\mathbb{E}_{\omega}(\omega|\omega \in \Omega^{l}))$$

#### Lower Bounding-Based LShaped Method

**0.** Let 
$$l = 1, \Omega_l = \Omega, k = 1, x^k = 0$$

**1.** For j = 1, 2, ... l solve...

$$\pi_j = \arg\max_{\pi \in \Re^m} \{ \pi^T (h(\bar{\omega}_j) - T(\bar{\omega}_j) x^k) | W^T \pi \le q \}$$

$$\bar{\omega}_j = \mathbb{E}_P(\omega | \omega \in \Omega_j)$$

2. Create cut as...

$$\theta \ge \sum_{j=1}^{l} \pi_j^T \left[ (h(\bar{\omega}_j) - T((\bar{\omega}_j)x^k)) \right] - \sum_{j=1}^{l} \pi_j^T \left[ \pi_j^T T((\bar{\omega}_j)) \right] (x - x_k)$$
$$(p_j = \mathsf{P}(\omega \in \Omega_j).$$

## Continuing

**3.** Let k = k + 1, add cut to master problem and solve

$$(x_k, \theta_k) = \arg\min_{x \in X} c^T x + \theta$$

subject to

$$\theta \ge \alpha_k + \beta_k x \quad \forall k = 1, 2, \dots$$

4.  $\mathcal{L}(x_k) = c^T x_k + \theta_x \leq \mathcal{Q}(x)$ . If you are happy with  $x_k$  based on  $\mathcal{L}(x_k)$ , then quit, otherwise create a new partition of  $\Omega = \{\Omega_1, \Omega_2, \dots, \Omega_l\}$ . Go to 1.

(I'll draw some pictures in class)...

**0.** Let 
$$k = 1$$
,  $x_k = 0$ ,  $V = \emptyset$ 

**1a.** Draw random sample  $\omega_k$ , and solve...

$$\pi_k = \arg \max_{\pi \in \Re^m} \{ \pi^T (h(\bar{\omega}_k) - T((\bar{\omega}_k)x^k) | W^T \pi \le q \}$$
  
1b.  $V = V \cup \pi^k$ . For  $j = 1, 2, \dots k - 1$ , solve

$$\pi^{j} = \arg\max_{\pi \in V} \left\{ \pi^{T} \left( (h(\bar{\omega}_{j}) - T((\bar{\omega}_{j})x^{k}) \right) \right\}$$

2a. Create cut as...

$$\theta \ge 1/k \sum_{j=1}^k \pi_j^T(h(\omega_j) - T(\omega_j)x_k)$$

• Call the cut  $(\alpha_k + \beta_k^T x)$ .

**2b.** For  $j = 1, 2, \ldots, k - 1$ , *Phase Out* old cuts as

$$\alpha_k + \beta_k^T x = (k-1)/k(\alpha_{k-1} + \beta_{k-1}^T x).$$

#### 3. Solve Master Problem

$$(x_k, \theta_k) = \arg\min_{x \in X} c^T x + \theta$$

subject to

$$\theta \ge \alpha_k + \beta_k x \quad \forall k = 1, 2, \dots$$

- Go to 1.
- $\star$  There is some *subsequence* of the  $x^k \to x^*$
- Typically people use some sort of statistical based stopping criteria

# LP Duality

• By duality

$$Q(x,\omega) = \max_{\pi \in \Re^m} \{\pi^T (h(\omega) - T(\omega)x) | \pi^T W \le q$$

• So if  $\hat{\pi}$  satisfies  $\hat{\pi}^T W \leq q$ , then

$$Q(x,\omega) \ge \hat{\pi}^T (h(\omega) - T(\omega)x) \quad \forall \omega \in \Omega, \forall x \in X$$

? (Why?)

⇒ The lower bounding function  $\mathcal{L}(x)$  is (only) such that  $\mathbb{E}_P \mathcal{L}(x) \leq \mathcal{Q}(x)$ 

#### **Checkpoint! – What we've learned**

- Modeling the Deterministic Equivalent of Stochastic (Linear) Programs
  - ◊ Two-stage
  - Multi-stage. (Modeling our favorite eight-syllable word).
- EVPI
- VSS
- Recourse Function

#### We Really Learned This Much?

- Two-Stage Stochastic LP With Recourse
  - Properties of the recourse function
    - Convex
    - Subdifferentiable
    - Lipschitz-Continuous
  - ◊ LShaped Method
    - Feasibility Cuts
    - Multicut methods
    - Regularized methods. Trust region and regularized decomposition.
    - Bunching and Trickling Down



- ♦ Bounds
  - Jensen's Inequality
  - Edmunson-Madansky Inequality
- Monte Carlo Methods
  - Valid Statistical Lower and Upper Bounds on optimal objective function value
  - ◊ Variance Reduction (Latin Hypercube Sampling)
  - Convergence of Optimal Solutions

#### Math Stuff We've Learned

- KKT Conditions
- Convexity
- (Lipschitz)-Continuity
- Minkowski Sum
- Weak-Strong Long of Large Numbers
- Central Limit Theorem
- Lebesgue-Stiltjes Integral
- L'Hopital's Rule

#### **Comouter Stuff We've Learned**

- AMPL! Lots and lots of AMPL.
- Numerical Integration
- MPS format
- SMPS
- High Performance Computing
  - ♦ SMP
  - Message Passing
- Grid Computing
  - ♦ Condor

## Survey Time

- A simple show of hands please...
- What Next?
  - ♦ Stochastic MIP
  - Probabilistic Constraints?
  - ♦ Multi-Stage Stochastic LP?
- Next Time...
  - Go over old homeworks
  - ♦ Get Homework "#5–#6"

#### **Remainder of Course**

- Start calling for people's projects to be done *soon!*
- Only 8 more classes left. Last 3–4 classes will be (interspersed) with project presentations ⇒ you have roughly 2–3 weeks to finish.
- There will be 1.5 more assignments
- $\Rightarrow$  There is *LOTS* of work that you have left to do for this course!