

IE 495 – Lecture 21

Multistage Stochastic Programming

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Outline

- HW Fixes
- Multistage Stochastic Programming
 - ◇ Modeling
 - ◇ Nested Decomposition Method

Homework Fixes

- Problem 2...
 - ◇ $\Omega^I = \{\omega_1 \times \omega_2 | 1 \leq \omega_1 \leq 5/2, 1/3 \leq \omega_2 \leq 2/3\}$
 - ◇ $\Omega^{II} = \{\omega_1 \times \omega_2 | 5/2 \leq \omega_1 \leq 4, 1/3 \leq \omega_2 \leq 2/3\}$
- Problem 3...
 - ◇ The constraints with ω should be *equality* constraints
- Problem 4...
 - ◇ Should be a *maximization* problem

Grade Data Collection

Stochastic Integer Programming (Please Don't Call On Me!)

- How would you solve a really small instance?
- How would you solve an instance with integer variables only in the first stage?
- What is an “optimality cut” in the context of the integer L-Shaped Method?
- For what problem are the optimality cuts we showed last time valid?
- Name one manner in which we might obtain lower and upper bounds to use in the stochastic branch and bound method

Jacob and MIT

- We are given a universe N of investment decisions
- We have a set $\mathcal{T} = \{1, 2, \dots, T\}$ of investment periods
- Let $\omega_{it}, i \in N, t \in \mathcal{T}$ be the return of investment $i \in N$ in period $t \in \mathcal{T}$.
- If we exceed our goal G , we get an interest rate of q that Helen and I can enjoy in our golden years
- If we don't meet the goal of G , Helen and I will have to borrow money at a rate of r so that Jacob can go to MIT.
- We have $\$b$ now.

Variables

- $x_{it}, i \in N, t \in \mathcal{T}$: Amount of money to invest in vehicle i during period t
- y : Excess money at the end of horizon
- w : Shortage in money at the end of the horizon

(Deterministic) Formulation

maximize

$$qy + rw$$

subject to

$$\sum_{i \in N} x_{i1} = b$$

$$\sum_{i \in N} \omega_{it} x_{i,t-1} = \sum_{i \in N} x_{it} \quad \forall t \in \mathcal{T} \setminus 1$$

$$\sum_{i \in N} \omega_{iT} x_{iT} - y + w = G$$

$$x_{it} \geq 0 \quad \forall i \in N, t \in \mathcal{T}$$

$$y, w \geq 0$$

One Way to Model

- One way to model this is to create copies of the variables for every scenario at every time period.
- Then we need to enforce *nonanticipativity*...
- Define S_s^t as the set of scenarios that are equivalent (or indistinguishable) to scenario s at time t

A Stochastic Version – Explicit Nonanticipativity

maximize

$$\sum_{s \in S} p_s (qy_s - rw_s)$$

subject to

$$\sum_{i \in N} x_{i1} = b$$

$$\sum_{i \in N} \omega_{its} x_{i,t-1,s} = \sum_{i \in N} x_{its} \quad \forall t \in \mathcal{T} \setminus 1, \forall s \in S$$

$$\sum_{i \in N} \omega_{iT} x_{iT_s} - y_s + w_s = G \quad \forall s \in S$$

$$x_{its} = x_{its'} \quad \forall i \in N, \forall t \in \mathcal{T}, \forall s \in S, \forall s' \in S_s^t$$

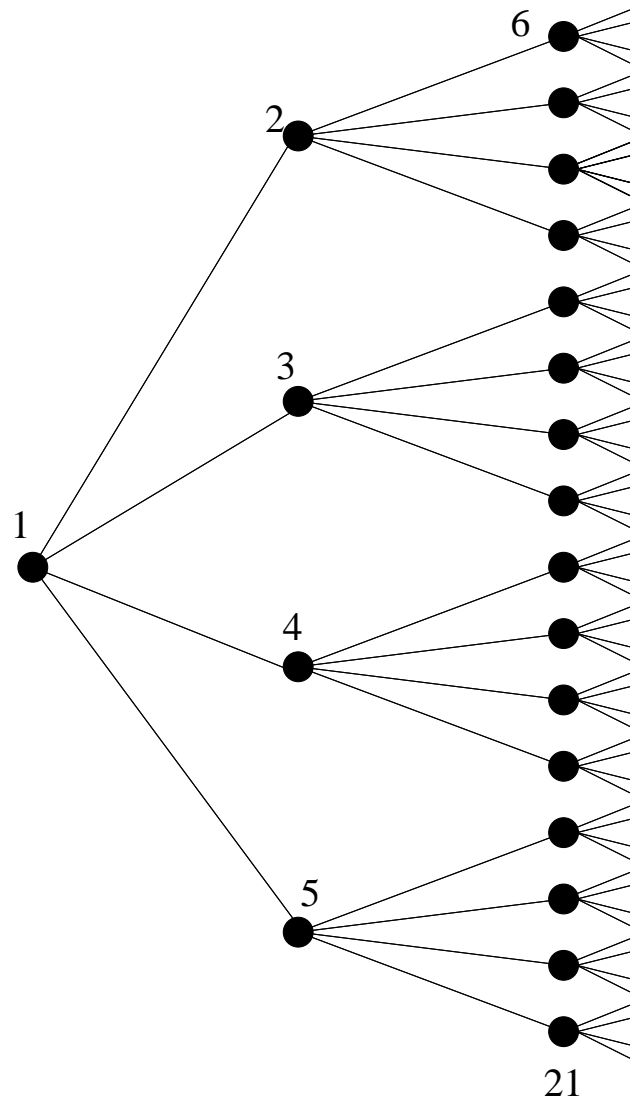
$$x_{its} \geq 0 \quad \forall i \in N, t \in \mathcal{T}, \forall s \in S$$

$$y_s, w_s \geq 0 \quad \forall s \in S$$

Another Way

- We can also enforce nonanticipativity by just not creating the “wrong” variables
- We have a vector of variables for each node in the tree.
- This vector corresponds to what our decision would be given the realizations of the random variables we have seen so far.
- Index the nodes $l = 1, 2, \dots, \mathcal{L}$.
- We will need to know the “parent” of any node.
- Let $A(l)$ be the ancestor of node $l \in \mathcal{L}$ in the scenario tree.

Jacob-MIT Event Tree



Another Multistage formulation

maximize

$$\sum_{s \in S} p_s (qy_s - rw_s)$$

subject to

$$\sum_{i \in N} x_{i1} = b$$

$$\sum_{i \in N} \omega_{il} x_{i,A(l)} = \sum_{i \in N} x_{il} \quad \forall l \in \mathcal{L} \setminus 1$$

$$\sum_{i \in N} \omega_{iA(s)} x_{iA(s)} - y_s + w_s = G \quad \forall s \in S$$

$$x_{il} \geq 0 \quad \forall i \in N, \forall l \in \mathcal{L}$$

$$y_s, w_s \geq 0 \quad \forall s \in S$$

Multistage Stochastic LP—Implicit Nonanticipativity

$$\min_{x_1} \left\{ c_1 x_1 + \mathbb{E}_{\omega_2} \left[c_2 x_2 + \mathbb{E}_{\omega_3 | \omega_2} \left[\min_{x_3} + \cdots + \mathbb{E}_{\omega_T | \omega_2, \dots, \omega_{T-1}} \left[\min_{x_T} c_T x_T \right] \right] \right] \right\}$$

$$\begin{array}{rcl}
 A_1 x_1 & & = h_1 \\
 A_2(\omega_2)x_1 + W_2 x_2(\omega_2) & & = h_2(\omega_2) \\
 & A(\omega_3)x_2(\omega_2) + W_3 x_3(\omega_3) & = h_3(\omega_3) \\
 & \vdots & \vdots \\
 & & A(\omega_T)x_{T-1}(\omega_{T-1}) + W_T x_T(\omega_T) = h_T(\omega_T) \\
 x_1 & & \geq 0 \\
 & x_2(\omega_2) & \geq 0 \\
 & \vdots & \geq 0 \\
 & & x_T(\omega_T) \geq 0
 \end{array}$$

Nested Decomposition Procedure

- A (The?) method for solving multistage stochastic programs
- Just like a recursive version of the L-Shaped Method
- Parent nodes send “proposals” for solutions to their children nodes
- Child nodes send cuts to their parents
- There are different “sequence procedures” that tell in which order the problems corresponding to different nodes in the scenario tree are solved.

Nested Decomposition

- “Fast Forward-Fast Back”
 - ◇ Most common sequencing procedure
 - ◇ Do all nodes at a time period
 - ◇ If you find a problem infeasible, create the feasibility cut and reverse direction
 - ◇ If problem has (relatively) complete recourse, it amounts to breadth-first search of on the scenario tree

Nested Decomposition

- Cuts have the form $(\theta \geq e - Ex)$, where

$$e = \sum_{k \in \mathcal{D}} (p_k | p_j) [\pi_k^T h_k + \sigma_k^T e_k]$$

$$E = \sum_{k \in \mathcal{D}} (p_k | p_j) \pi_k^T T_k$$

Hot! Hot! Hot! From B&L Section 7.1

- Planning Production of Air Conditioners (ACs) for next three months
- In each month we can produce at most 200 ACs for \$100 each
- We can use overtime workers if the demand is heavy, but then the cost per unit is \$300
- The demand in the first month is 100
- The demand in the second and third months is random. With 50% probability it will be either 100 or 300.
- We can store ACs at a cost of \$50/month
- We *must* meet all demand
- No salvage/disposal costs at the end of the horizon

Hot! Hot! Hot! Model

Variables

- x_t —Number of ACs to produce (regular) in time t
- w_t —Number of ACs to produce (overtime) in time t
- y_t —Number of ACs to carry over in inventory at the end of time t
- d_t —Demand for ACs in time t

Model at node l

$$\min x_t + 3w_t + 0.5y_t + \theta_l$$

subject to

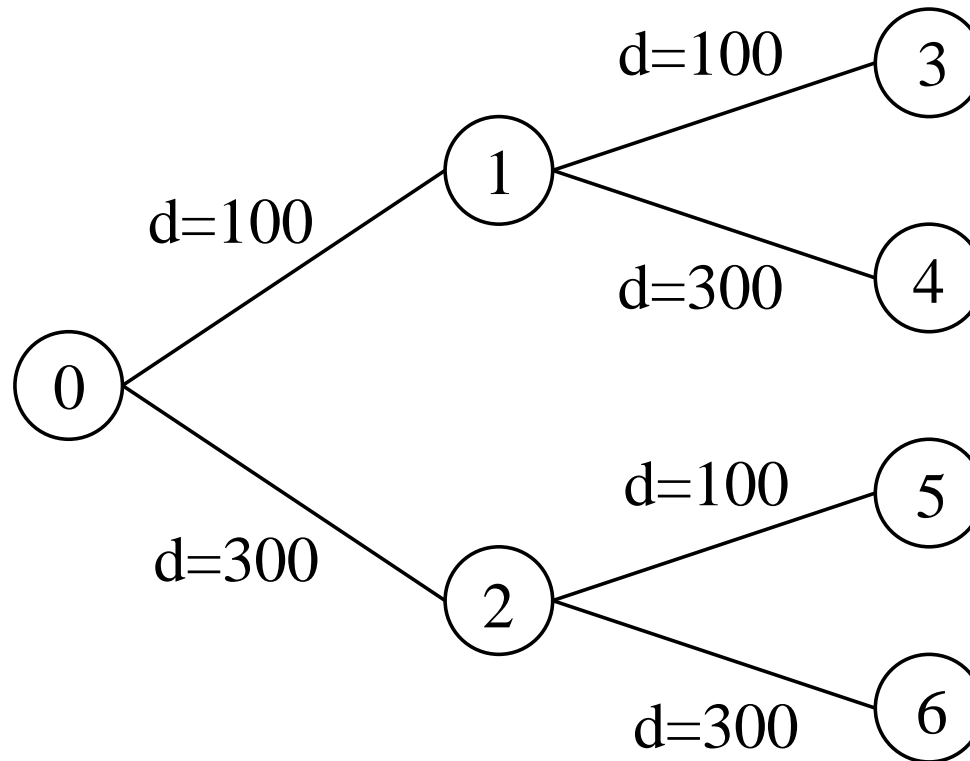
$$x_t \leq 2$$

$$y_{t-1} + x_t + w_t - y_t = d_l$$

$$\theta_l \geq e_k - E_k(x, w, y)^T \quad \forall \text{cuts at this node}$$

$$\text{All vars} \geq 0$$

Scenario Tree



- We're going to do a couple iterations of Nested Decomposition.
- **Pay Attention**, it will probably be on the final

April 23 Presentations...

(Only going to do five people, since you will have to fill out evaluations on that day too...)

- Banu Gemicı
 - Rui Kang
 - Jen Rogers
 - Jerry Shen
 - Clara Novoa
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- Next time—Probabilistic Constraints