

IE 495 – Lecture 8

SMPS and the Recourse Function

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Outline

- Formulating Stochastic Program(s).
 - ◇ SMPS format
 - ◇ Jacob & MIT
- Two stage problems with recourse – Expected value function

Please don't call on me!

- What are the KKT conditions?
- Who is Karmarkar?
- The simplex method is a polynomial time algorithm for linear programming.
- True or False: MPS format is a concise, new format for expressing linear and integer programs?
- Explain two ways to model “your favorite eight letter word”?

Multistage Formulation—Implicit Nonanticipativity

maximize

$$\sum_{s \in S} p_s (qy_s - rw_s)$$

subject to

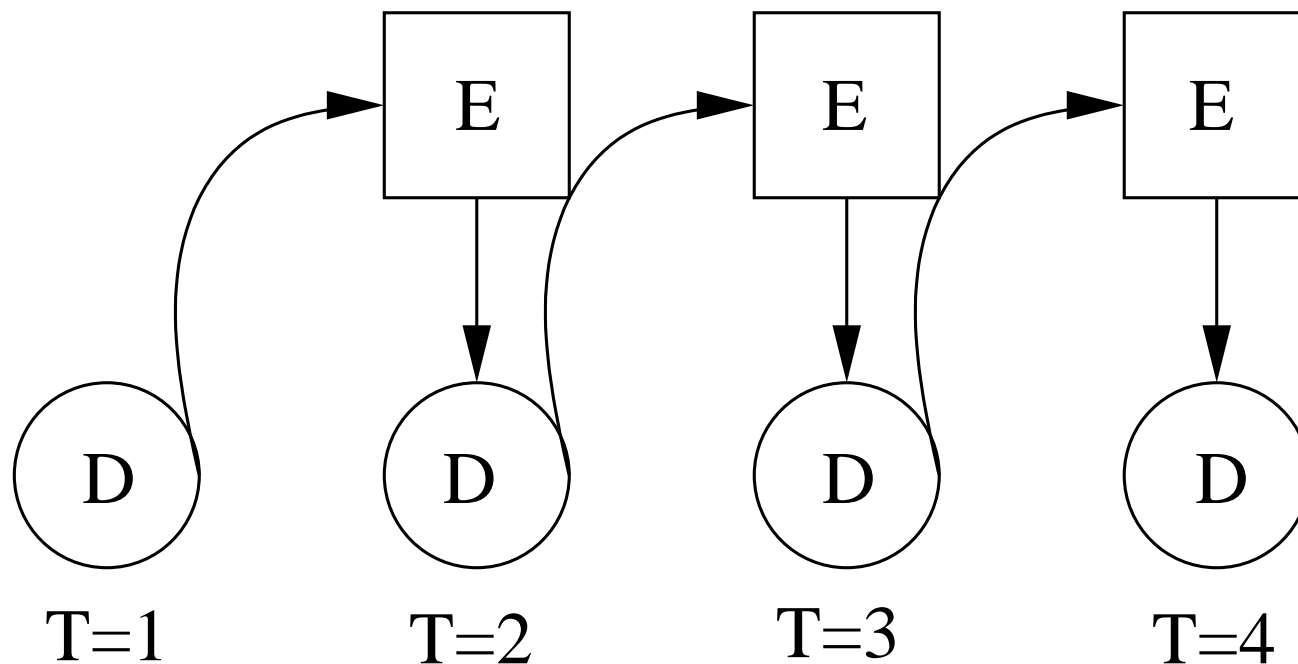
$$\sum_{i \in N} x_{i1} = b$$

$$\sum_{i \in N} \omega_{il} x_{i,A(l)} = \sum_{i \in N} x_{il} \quad \forall l \in \mathcal{L} \setminus 1$$

$$\sum_{i \in N} \omega_{iA(s)} x_{iA(s)} - y_s + w_s = G \quad \forall s \in S$$

SMPS

- Multistage problems are based on the “event-decision” model



- All random “stuff” must be in the stage associated with the decision.

ω in “Wrong” Stage

- To get SMPS to “work”. Let’s rewrite the problem so that the ω_{it} are associated with variables in stage t
- Let x_{it} : Amount of money invested in i in time t
- Let y_{it} : Amount of money you have in i at time t
- $(1/\omega_{it})y_{it} = x_{i,t-1}$

An Equivalent (Longer) Formulation

maximize

$$qE + rU$$

subject to

$$\sum_{i \in N} x_{i1} = b$$

$$\omega_{it}^{-1} y_{it} - x_{i,t-1} = 0 \quad \forall i \in N, \forall t \in \mathcal{T} \setminus 1$$

$$\sum_{i \in N} x_{it} - \sum_{i \in N} y_{it} = 0 \quad \forall t \in \mathcal{T} \setminus 1$$

$$\sum_{i \in N} y_{iT} - E + B = G$$

$$x_{it} \geq 0 \quad \forall i \in N, t \in \mathcal{T}$$

$$y_{it} \geq 0 \quad \forall i \in N, t \in \mathcal{T} \setminus 1$$

$$E, U \geq 0$$

The Matrix

- $N = \{S, B\}$
- $T=4$

x_{S1}	x_{B1}	y_{S2}	y_{B2}	x_{S2}	x_{B2}	y_{S3}	y_{B3}	x_{S3}	x_{B3}	y_{S4}	y_{B4}	E	U	
1	1													b
-1		$1/\omega_S$												0
	-1		$1/\omega_B$											0
		-1	-1	1	1									0
				-1		$1/\omega_S$								0
					-1		$1/\omega_B$							0
						-1	-1	1	1					0
								-1		$1/\omega_S$				0
									-1		$1/\omega_B$			0
										1	1	-1	1	G

SMPS Format

- How do we specify a stochastic programming instance to the solver?
- We could form the deterministic equivalent ourselves, but you saw how unnatural that seems.
 - ◇ For *really* big problems, forming the deterministic equivalent is out of the questions.
 - ◇ We need to just specify the random parts of the model.
- We can do this using SMPS format
 - ◇ Actually, other than (explicitly) forming the deterministic equivalent in a modeling language, this is the **ONLY** way.
 - ★ There *is* some recent research work in developed stochastic programming support in an AML.

SMPS Components

- Core file
 - ◇ Like MPS file for “base” instance
- Time file
 - ◇ Specifies the time dependence structure
- Stoch file
 - ◇ Specifies the randomness

SMPS

- The SMPS format is *broad*^a
 - ◇ There are very few (if any) full implementations, that simply read the full format
 - ◇ No solver will solve all instances that can be expressed in the format.
- The SMPS format is (seemingly) being changed.
- A good site...
 - ◇ <http://www.mgmt.dal.ca/sba/profs/hgassmann/SMPS2.htm>

^a *Too* broad, IMO

Papers

- I handed out to you...
 - ◇ J.R. Birge, M.A.H. Dempster, H.I. Gassmann, E.A. Gunn, A.J. King and S.W. Wallace, “A standard input format for multiperiod stochastic linear programs”, COAL Newsletter #17 (1987) pp. 1-19.
- There are more. For example...
 - ◇ H.I. Gassmann and E. Schweitzer, “A comprehensive input format for stochastic linear programs”, Annals of Operations Research 104 (2001) 89-125.
- Any questions?

SMPS Core File

- Like an MPS file specifying a “base” scenario
- ★ *Must* permute the rows and columns so that the time indexing is sequential

```

NAME      jake
ROWS
N  obj
E  c1
E  c2
E  c3
E  c4
E  c5
E  c6
E  c7
E  c8
E  c9
E  c10
COLUMNS
    xs1      c1          1  c2          -1
    xb1      c1          1  c3          -1
    ys2      c2         999  c4          -1
    yb2      c3         888  c4          -1
    xs2      c4          1  c5          -1
    xb2      c4          1  c6          -1
    ys3      c5         999  c7          -1
    yb3      c6         888  c7          -1
    xs3      c7          1  c8          -1
    xb3      c7          1  c9          -1
    ys4      c8         999  c10         1
    yb4      c9         888  c10         1
    t        obj        -0.05 c10         -1
    s        obj         1.1  c10         1
RHS
    rhs      c1         10000  c10         15000
ENDATA

```

jake.time

- ◇ Specify which row/column starts each time period.
- ★ Must be sequential!

```
*23456789 123456789 123456789
```

```
TIME           jake
```

```
PERIODS       IMPLICIT
```

```
    xs1        c1           T1
```

```
    ys2        c2           T2
```

```
    ys3        c5           T3
```

```
    ys4        c8           T4
```

```
ENDATA
```

Stoch File

- BLOCKS
 - ◇ Specify a “block” of parameters that changes together
- INDEP
 - ◇ Specify that all the parameters you are specifying are all independent random variables
- SCENARIO
 - ◇ Specify a “base” scenario
 - ◇ Specify what things change and when...

jake.stoch

```
*23456789 123456789
STOCH      jake
*23456789 123456789 123456789 123456789
BLOCKS     DISCRETE
BL BLOCK1  T2      0.5
            ys2     0.8
            yb2     0.8772
BL BLOCK1  T2      0.5
            ys2     0.9434
            yb2     0.8929
BL BLOCK2  T3      0.5
            ys3     0.8
            yb3     0.8772
BL BLOCK2  T3      0.5
            ys3     0.9434
            yb3     0.8929
BL BLOCK3  T4      0.5
            ys4     0.8
            yb4     0.8772
BL BLOCK3  T4      0.5
            ys4     0.9434
            yb4     0.8929
ENDATA
```

An INDEP Example

*23456789 123456789

STOCH jake

*23456789 123456789 123456789 123456789

INDEP DISCRETE

ys2	c2	0.8	0.333
-----	----	-----	-------

ys2	c2	1.0	0.333
-----	----	-----	-------

ys2	c2	1.05	0.333
-----	----	------	-------

*

yb2	c3	0.9	0.125
-----	----	-----	-------

yb2	c3	0.99	0.675
-----	----	------	-------

yb2	c3	1.0	0.125
-----	----	-----	-------

ENDATA

NEOS

- If we have files in SMPS format, how do we solve the resulting instance?
- We can use NEOS!

<http://www.mcs.anl.gov/neos>

- There are currently two solvers
 - ◇ CPA—Works for two-stage LP
 - ◇ MSLIP—Works for multistage LP
- I am working on getting others that we can use.
- ★ You will (likely) be asked to solve some an instance on your next homework assignment

Math Time! Two-Stage SLP w/Fixed Recourse

minimize

$$c^T x + \mathbb{E}_\omega [q^T y]$$

subject to

$$Ax = b$$

$$T(\omega)x + Wy(\omega) = h(\omega) \quad \forall \omega \in \Omega$$

$$x \in X$$

$$y(\omega) \in Y$$

-
- $Q(x, \omega) = \min_{y \in Y} \{q^T y : Wy = h(\omega) - T(\omega)x\}$

Some Notation Review

$$\min_{x \in X: Ax=b} \left\{ c^T x + \mathbb{E}_\omega \left[\min_{y \in Y} \{ q^T y : Wy = h(\omega) - T(\omega)x \} \right] \right\}$$

- *Second stage value function, or recourse (penalty) function*
 $v : \Re^m \mapsto \Re$.
- $v(z) \equiv \min_{y \in Y} \{ q^T y : Wy = z \}$,
 - ◇ Given “policy” x and realization of randomness ω
 - ◇ If z measures the first-stage deviation $z = h(\omega) - T(\omega)x$,
 $v(z)$ is the minimum cost way to “correct” so that the constraints hold again.
- $Q(x, \omega) = v(h(\omega) - T(\omega)x)$

More Notation

- *Expected Value Function, or Expected minimum recourse function* $Q : \mathcal{R}^n \mapsto \mathcal{R}$.
 - ◇ $Q(x) \equiv \mathbb{E}_\omega[Q(x, \omega)]$
 - ◇ For any policy $x \in \mathcal{R}^n$, it describes the expected cost of the recourse.

The SP Problem

- Using these definitions, we can write our recourse problem in terms only of the x variables:

$$\min_{x \in X} \{c^T x + Q(x) : Ax = b\}$$

- This is a (nonlinear) programming problem in \mathcal{R}^n .
- ⇒ The ease of solving such a problem depends on the properties of $Q(x)$.
- ? Does anyone know what $Q(x)$ is?
 - ◇ Linear,(?) Convex,(?) Continuous,(?) Differentiable(?)
- This is what we are going to study for a while...

Here We Go

- For the time being, let $Y = \mathfrak{R}_+^p$.

$$v(z) = \min_{y \in \mathfrak{R}_+^p} \{q^T y : Wy = z\}, z \in \mathfrak{R}^m$$

- Thus, for a fixed z , we solve a linear program to evaluate $v(z)$.
- Assume that LP duality holds $\forall z \in \mathfrak{R}^m$
 - ◇ Only for the time being!
 - ◇ We'll reconsider later...
 - ◇ $-\infty < v(z) < \infty$
- ★ This is some notation. Let
 - $\{y \in \mathfrak{R}_+^p : Wy = z\} = \emptyset \Rightarrow v(z) = \infty$
 - $\exists d \in \mathfrak{R}_+^n : Wd = 0, q^T d < 0 \Rightarrow v(z) = -\infty$.

Proofs

- So if LP duality holds...

$$v(z) = \min_{y \in \mathfrak{R}_+^p} \{q^T y : W y = z\} = \max_{t \in \mathfrak{R}^m} \{z^T t : W^T t \leq q\}$$

- Let $\Lambda = \{\lambda_1, \lambda_2, \dots, \lambda_{|\Lambda|}\}$ be the set of extreme points of $\{t \in \mathfrak{R}^m \mid W^T t \leq q\}$.
 - ◇ Each of those extreme points λ_k is potentially an optimal solution to the LP.
 - ◇ In fact, we are sure that there is no optimal solution better than one that occurs at an extreme point, so we can write...

$$v(z) = \max_{k=1, \dots, |\Lambda|} \{z^T \lambda_k\}, z \in \mathfrak{R}^m.$$

What's All This?

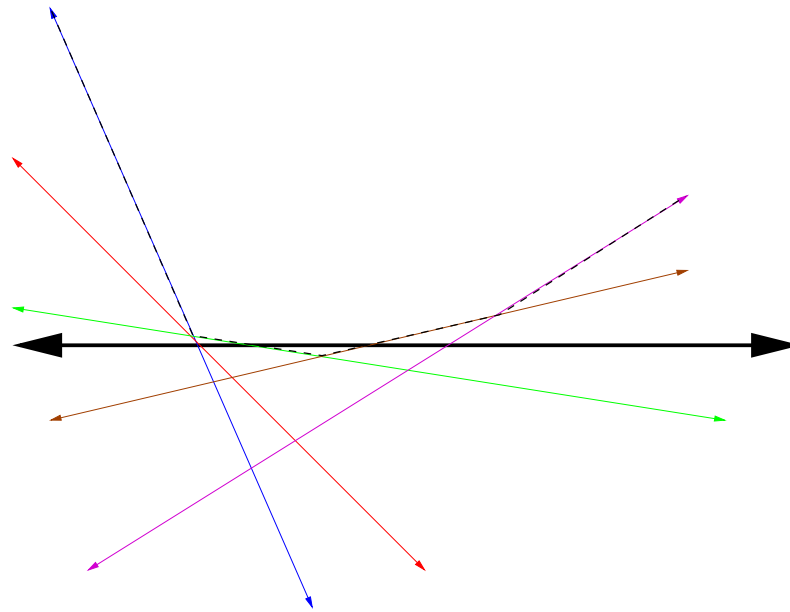
$$\begin{aligned}\alpha v(z_1) + (1 - \alpha)v(z_2) &= \max_{k=1,2,\dots,|\Lambda|} z_1^T \lambda_k + (1 - \alpha) \max_{k=1,2,\dots,|\Lambda|} z_2^T \lambda_k \\ &\leq \alpha z_1^T \lambda_k^* + (1 - \alpha) z_2^T \lambda_k^* \\ &= (\alpha z_1 + (1 - \alpha)z_2)^T \lambda_k^* \\ &\leq \max_{k=1,2,\dots,|\Lambda|} [(\alpha z_1 + (1 - \alpha)z_2)^T \lambda_k] \\ &= v((\alpha z_1 + (1 - \alpha)z_2))\end{aligned}$$

Quite Enough Done

? What did I just prove?

Convex!

- $v(z)$ is convex of $z \in \mathfrak{R}^m$.
- In fact...
 - ◇ **Thm:** If $f_1(x), f_2(x), \dots, f_q(x)$ is an arbitrary collection of convex functions, then $M(x) = \max\{f_1(x), f_2(x), \dots, f_q(x)\}$ is also a convex function.



I Don't Care About v

? What about $Q(x, \omega)$?

◇ Recall $Q(x, \omega) \equiv v(h(\omega) - T(\omega)x)$

$$\begin{aligned} & \lambda Q(x_1, \omega) + (1 - \lambda)Q(x_2, \omega) \\ &= \lambda v(h(\omega) - T(\omega)x_1) + (1 - \lambda)v(h(\omega) - T(\omega)x_2) \\ &\geq v(\lambda(h(\omega) - T(\omega)x_1) + (1 - \lambda)(h(\omega) - T(\omega)x_2)) \\ &= v(h(\omega) - T(\omega)(\lambda x_1 + (1 - \lambda)x_2)) \\ &= Q(\lambda x_1 + (1 - \lambda)x_2, \omega) \end{aligned}$$

Quite Enough Done

Continuing On

- So $Q(x, \omega)$ is convex in x for a fixed ω .
- In fact...
 - ◇ **Thm:** If A is a linear transformation from $\mathfrak{R}^n \mapsto \mathfrak{R}^m$, and $f(x)$ is a convex function on \mathfrak{R}^m , the composite function $(fA)(x) \equiv f(Ax)$ is a convex function on \mathfrak{R}^n .

Almost Done...

- What about $Q(x) \equiv \mathbb{E}_\omega Q(x, \omega)$
 - ◇ From the remainder of today, let's assume that ω comes from a probability space with finite support.
 - ◇ This means that there are finite number of discrete values $\{\omega_1, \omega_2, \dots, \omega_m\}$ that ω can take.

$$Q(x) = \sum_{i=1}^m P(\omega = \omega_i) Q(x, \omega_i)$$

Finishing up

- **Thm:** If $f(x)$ is convex, and $\alpha \geq 0$, $g(x) \equiv \alpha f(x)$ is convex.
- **Thm:** If $f_k(x)$, $k = 1, 2, \dots, K$ are convex functions, so is $g(x) \equiv \sum_{k=1}^K f_k(x)$.
- Put it all together and you get...
 - ★ $Q(x)$ is a convex function of x .

What assumptions have we made so far?

- LP duality holds $\forall z = (h(\omega) - T(\omega)x) \in \mathfrak{R}^m$
- ω has finite support

Next time

- Discuss these assumptions
- Show more properties of $Q(x)$
 - ◇ Optimality conditions
- Homework #2. :-(
 - ◇ It will be less time consuming than HW#1